

Segment Trees

CS 491 – Competitive Programming

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Running Example

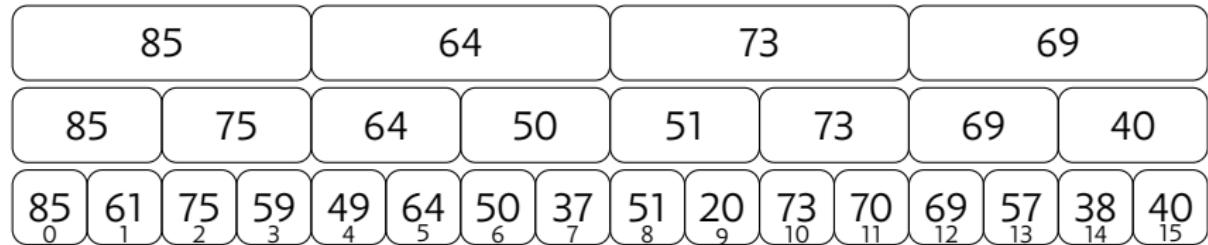
- ▶ Consider the following array:

85 0	61 1	75 2	59 3	49 4	64 5	50 6	37 7	51 8	20 9	73 10	70 11	69 12	57 13	38 14	40 15
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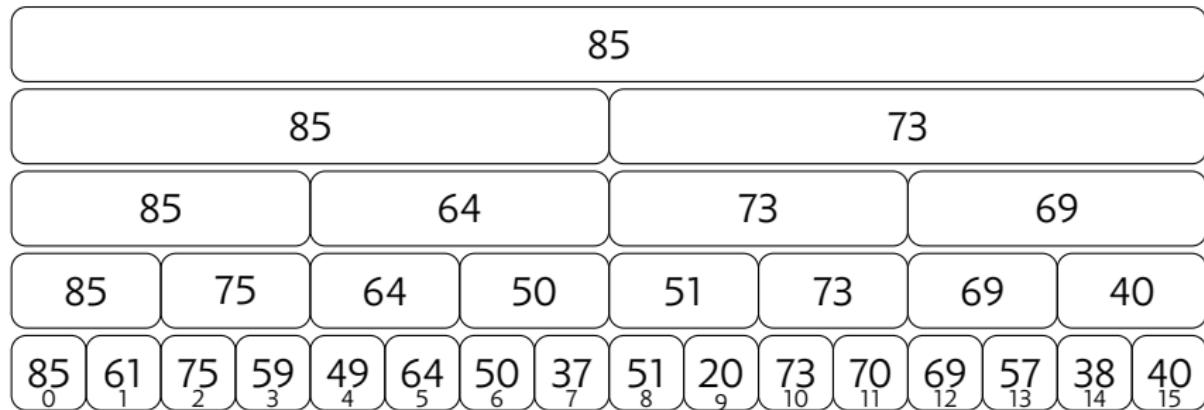
Segment Trees, Level 1

85	75	64	50	51	73	69	40								
85 ₀	61 ₁	75 ₂	59 ₃	49 ₄	64 ₅	50 ₆	37 ₇	51 ₈	20 ₉	73 ₁₀	70 ₁₁	69 ₁₂	57 ₁₃	38 ₁₄	40 ₁₅

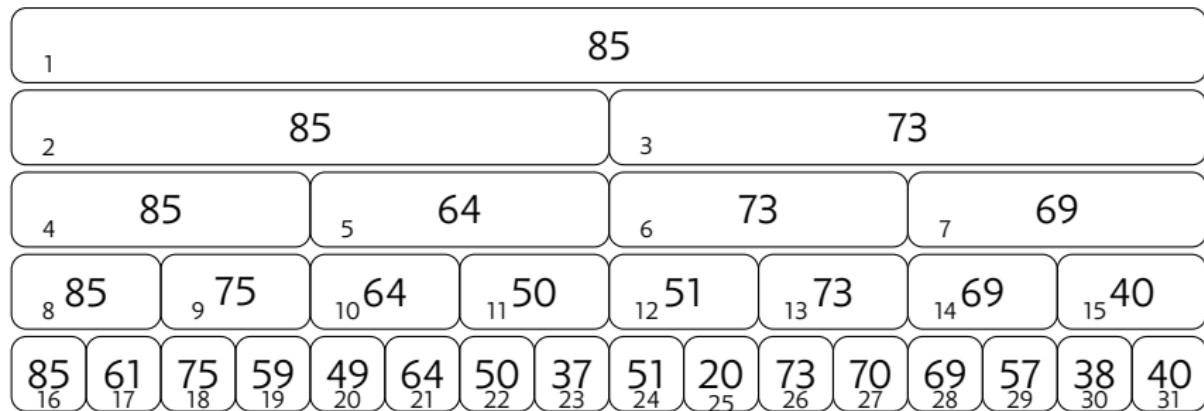
Segment Trees, Level 2



Segment Trees, Level 3 and 4



Segment Trees, Numering the Elements



Build the Segement Tree

- ▶ L and R give you the bounds with respect to the original array.
- ▶ This code gives you a max range.

```
1 void build(int p, int L, int R) {  
2     if (L == R) // as L == R, either one is fine  
3         st[p] = data[L]; // store the data  
4     else {  
5         // recursively compute the values  
6         build(left(p) , L           , (L + R) / 2);  
7         build(right(p), (L + R) / 2 + 1, R           );  
8         int p1 = st[left(p)], p2 = st[right(p)];  
9         st[p] = (st[p1] >= st[p2]) ? p1 : p2;  
10    } }
```

Query the Tree

- ▶ L and R give you the bounds with respect to the original array.
- ▶ i and j give you the bounds for the query

```
1 int rmq(int p, int L, int R, int i, int j) {  
2     if (i > R || j < L) return -1; // current segment outside query  
3     if (L >= i && R <= j) return st[p];  
4     // compute the min position in the left and right part of the segment  
5     int p1 = rmq(left(p), L, (L+R)/2, i, j);  
6     int p2 = rmq(right(p), (L+R)/2+1, R, i, j);  
7     if (p1 == -1) return p2;  
8     // if we try to access segment outside query  
9     if (p2 == -1) return p1;  
10    return (p1 >= p2) ? p1 : p2;  
11}  
12
```